## Lab 3 and 4

## Contents

1 Highlighting vertices in a graph 1
2 Part 1: Test for a vertex cover 2
3 Part 2: Maximal independent set 3
3.1 Test for independence . . . . . . . . . . . . . . . . . . . . . . 3
3.2 Maximal independent set . . . . . . . . . . . . . . . . . . . . 3

## 1 Highlighting vertices in a graph

When plotting a graph, it is possible to color certain vertices with another color. This can be achieved with highlight function as presented in the following example:

```
clear; close all;
A = [
    0 1 0 1;
    1 0 1 0;
    0 1 0 1;
    10 1 0
    ]; % Adjacency matrix
C = [1, 3]; % Vertices to highlight
figure;
G = graph(A);
h = plot(G);
highlight(h, C, 'NodeColor', 'r'); % 'r' for red, 'g'
    for green, 'b' for blue, etc.
```


## 2 Part 1: Test for a vertex cover

Write a program that for a given undirected graph $G=(V, E)$ and a subset of vertices $C$ of $V$ checks whether $C$ is a vertex cover of $G$. Plot the graph and highlight the vertices of $C$ if $C$ is a vertex cover.

Example of an input is:

```
E = [
    1, 2;
    2, 3;
    4, 5;
    6, 7;
];
C = [1, 2, 3, 4, 5, 6, 7]; % Set of vertices (vertex
    cover)
```

You can use the following pseudocode:

```
Algorithm 1: Test for a vertex cover
    Input : Edge list \(E\), subset of vertices \(C\)
    Output: Boolean value is_vc (true if \(C\) is a vertex cover of \(G\), false
                    otherwise)
    \(i s \_v c=\) true;
    for every edge \(u, v\) in \(E\) do
        covered \(=\) false;
        for \(j=1 \mathbf{T O}|C|\) do
            if \(C[j]==u\) or \(C[j]==v\) then
                covered \(=\) true;
            end
        end
        if not covered then
            is_vc = false;
        end
    end
```


## 3 Part 2: Maximal independent set

### 3.1 Test for independence

Write a program which for a given subset of vertices $I$ of an undirected graph $G=(V, E)$ checks whether $I$ is an independent set of $G$. Plot the graph and highlight the vertices which are in the IS.

You can base your code on the following pseudocode that uses an adjacency matrix representation of $G$ :

```
Algorithm 2: Test for independence
    Input : Adjacency matrix \(A\), set of vertices \(I\)
    Output: Boolean value \(t\) (true if \(I\) is an independent set and false
                otherwise)
    \(t=\) true;
    if \(|I|>1\) then
        for \(v_{\text {index }}=1 \boldsymbol{T O}|I|-1\) do
            \(v=I\left[v_{\text {index }}\right]\);
            for \(w_{\text {index }}=v_{\text {index }}+1 \boldsymbol{T O}|I|\) do
                \(w=I\left[w_{\text {index }}\right] ;\)
                if there is an edge between \(v\) and \(w\) then
                    \(t=\) false;
            end
        end
        end
    end
```


### 3.2 Maximal independent set

Write a program which for a given graph $G=(V, E)$ finds a maximal Independent Set in it.

You can implement a simple greedy algorithm which is based on the idea of starting from an empty set $I$, picking vertices of $G$ one by one and adding them to $I$ if $I$ continues to be an independent set.

The algorithm is described in more details in the following pseudocode:

```
Algorithm 3: Maximal Independent Set
    Input : Adjacency matrix \(A\)
    Output: A subset of vertices \(I\)
    \(P=\) a random permutation of vertices \(V\);
    \(I=\varnothing ;\)
    for every vertex \(v\) in \(P\) do
        \(T=I \cup v ;\)
        if \(T\) is an Independent Set then
            \(I=I \cup v ;\)
        end
    end
```

Hint: To find a random permutation of vertices of $G$ you can use the following command: $\mathrm{P}=\operatorname{randperm}(\operatorname{size}(\mathrm{A}, 1))$;

